

# The degree of functions in the Johnson and $q$ -Johnson schemes

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joint work with Jonathan Mannaert and Alfred Wassermann

## Introductory remarks

- ▶ Joint work with  
Jonathan Mannaert and Alfred Wassermann.
- ▶ Despite title  
“The degree of functions  
in the Johnson and  $q$ -Johnson schemes”  
No association schemes in this talk!
- ▶ Motivation (next slide) is geometric.  
Indeed: Topic close to design theory.  
Studied objects are “dual designs”.

## Cameron-Liebler line classes

- ▶ Cameron, Liebler 1982:  
“Special” set  $\mathcal{L}$  of **lines** in  $\text{PG}(3, q)$ .
- ▶ Defined by the following equivalent properties:
  - ▶ Algebraic property:  
 $\chi_{\mathcal{L}} \in \mathbb{R}$ -row space of the **point-line** incidence matrix.
  - ▶ Geometric property:  
Constant intersection with any line spread of  $\text{PG}(3, q)$ .

## In literature: Various directions of generalization

- ▶ Ambient space  $\text{PG}(n, q)$ .
- ▶ **lines**  $\longrightarrow$   **$k$ -spaces**.
- ▶ Allow  $q = 1$  (set case).
- ▶ **points**  $\longrightarrow$  spaces of degree  **$t$** .

## Goal

Coherent theory of all above generalizations.

## Subset and subspace lattices

- ▶ Fix  $q = 1$  (set case) or prime power  $q \geq 2$  ( $q$ -analog case).
- ▶ Fix  $n$  non-negative integer.
- ▶ Let  $V$  be a  $\begin{cases} \text{set of size } n \\ \mathbb{F}_q\text{-vector space of dimension } n \end{cases}$
- ▶ Let  $\mathcal{L}(V)$  be the lattice of all  $\begin{cases} \text{subsets of } V \\ \mathbb{F}_q\text{-subspaces of } V \end{cases}$
- ▶ For  $U \in \mathcal{L}(V)$  let  $\text{rk}(U) = \begin{cases} \#U \\ \dim(U) \end{cases}$
- ▶ Let  $\begin{bmatrix} V \\ k \end{bmatrix} = \{U \in \mathcal{L}(V) \mid \text{rk}(U) = k\}$ .  
Set case:  $\# \begin{bmatrix} V \\ k \end{bmatrix} = \binom{n}{k} = \begin{bmatrix} n \\ k \end{bmatrix}_1$  Binomial coefficient.  
 $q$ -analog case:  $\# \begin{bmatrix} V \\ k \end{bmatrix} = \begin{bmatrix} n \\ k \end{bmatrix}_q$  Gaussian coefficient.
- ▶ Always: Use algebraic dimension!  
(Except in established symbols like  $\text{PG}(n, q)$ ).

## Algebraic property

- ▶ Algebraic property of Cameron-Liebler line classes:  
 $\chi_{\mathcal{L}} \in \mathbb{R}$ -row space of the point-line incidence matrix.
- ▶ Straightforward generalization:
  - ▶ Let  $W^{(tk)}$  incidence matrix of  $t$ -spaces vs.  $k$ -spaces.
  - ▶ Let  $V_t$  be the  $\mathbb{R}$ -row space of  $W^{(tk)}$ .
  - ▶ Function  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$  has algebraic property  $A_t$  if  $f \in V_t$ .

## Baby example

- ▶ Let  $q = 1$ ,  $V = \{1, 2, 3, 4, 5, 6\}$  (so  $n = 6$ ),  $k = 3$ ,  $t = 2$ .
- ▶ Let  $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\} \subseteq \begin{bmatrix} V \\ 3 \end{bmatrix}$ .
- ▶ Claim: Set  $\mathcal{F}$  has algebraic property  $A_2$ ,  
i. e. its characteristic function  $\chi_{\mathcal{F}} : \begin{bmatrix} V \\ 3 \end{bmatrix} \rightarrow \mathbb{R}$  has prop.  $A_2$ .

## Baby example (cont.)

$W^{(23)} =$

	123	124	125	126	134	135	136	145	146	156	234	235	236	245	246	256	345	346	356	456
12	1	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
13	1	0	0	0	1	1	1	0	0	0	0	0	0	0	0	0	0	0	0	0
14	0	1	0	0	1	0	0	1	1	0	0	0	0	0	0	0	0	0	0	0
15	0	0	1	0	0	1	0	1	0	1	0	0	0	0	0	0	0	0	0	0
16	0	0	0	1	0	0	1	0	1	1	0	0	0	0	0	0	0	0	0	0
23	1	0	0	0	0	0	0	0	0	0	1	1	1	0	0	0	0	0	0	0
24	0	1	0	0	0	0	0	0	0	0	1	0	0	1	1	0	0	0	0	0
25	0	0	1	0	0	0	0	0	0	0	0	1	0	1	0	1	0	0	0	0
26	0	0	0	1	0	0	0	0	0	0	0	1	0	1	1	0	0	0	0	0
34	0	0	0	0	1	0	0	0	0	0	1	0	0	0	0	0	1	1	0	0
35	0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	0	1	0	1	0
36	0	0	0	0	0	0	1	0	0	0	0	0	1	0	0	0	0	1	1	0
45	0	0	0	0	0	0	0	1	0	0	0	0	0	1	0	0	1	0	0	1
46	0	0	0	0	0	0	0	0	1	0	0	0	0	0	1	0	0	1	0	1
56	0	0	0	0	0	0	0	0	0	1	0	0	0	0	0	1	0	0	1	1

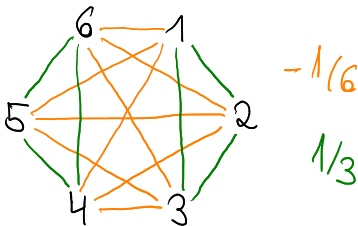
[illegible]

## Baby example (cont.)

- ▶  $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$ .
- ▶ We found:  $\mathcal{F}$  has property  $A_2$   
and the vector of **2-weights** of  $\mathcal{F}$  is

$$\text{wt}_{\mathcal{F}}^{(2)} = \left(\frac{1}{3}, \frac{1}{3}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, \frac{1}{3}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, -\frac{1}{6}, \frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right).$$

- ▶ Visualization.



- ▶ Exercise.

$\mathcal{F}$  does **not** have  $A_1$ .

## Geometric property

- ▶ Geometric property of Cameron-Liebler line classes:  
Constant intersection with any line spread of  $\text{PG}(3, q)$
- ▶ Generalization? – Not so clear.
- ▶ Observation:  
line spread of  $\text{PG}(3, q)$   
= set of lines in  $\text{PG}(3, q)$  covering every point exactly once  
= simple  $1-(4, 2, 1)_q$  subspace design
- ▶  $\rightsquigarrow$  use designs!

## Definition: Simple design

A set  $\mathcal{D} \subseteq \begin{bmatrix} V \\ k \end{bmatrix}$  is called a simple  $t-(n, k, \lambda)_q$  design, if every  $T \in \begin{bmatrix} V \\ t \end{bmatrix}$  is contained in exactly  $\lambda$  elements of  $\mathcal{D}$ .

- ▶ set case  $q = 1$ : combinatorial design
- ▶  $q$ -analog case  $q \geq 2$ : subspace design



## Example

► Let  $q = 1$ ,  $V = \{1, 2, 3, 4, 5, 6\}$  (so  $n = 6$ ),  $k = 3$ ,  $t = 2$ .

► Let

$$\mathcal{D} = \left\{ \{1, 2, 3\}, \{1, 2, 4\}, \{1, 3, 6\}, \{1, 4, 5\}, \{1, 5, 6\}, \right. \\ \left. \{2, 4, 6\}, \{2, 5, 6\}, \{2, 3, 5\}, \{3, 4, 5\}, \{3, 4, 6\} \right\} \subseteq \begin{bmatrix} V \\ 3 \end{bmatrix}.$$

► Check design condition for  $t = 2$ .

►  $T = \{1, 2\}$  is contained in blocks  $\{1, 2, 3\}$  and  $\{1, 2, 4\}$ .

►  $T = \{1, 3\}$  is contained in blocks  $\{1, 2, 3\}$  and  $\{1, 3, 6\}$ .

► ...

►  $T = \{5, 6\}$  is contained in blocks  $\{1, 5, 6\}$  and  $\{2, 5, 6\}$ .

►  $\implies \mathcal{D}$  is simple  $2$ -( $6, 3, 2$ )<sub>1</sub> design.

## Example (Trivial simple designs)

►  $\emptyset$  is empty  $t$ -( $v, k, 0$ )<sub>q</sub> design.

►  $\begin{bmatrix} V \\ k \end{bmatrix}$  is complete  $t$ -( $v, k, \lambda_{\max}$ )<sub>q</sub> design  
where  $\lambda_{\max} := \begin{bmatrix} n-t \\ k-t \end{bmatrix}$ .

## Definition: Simple design (repeated)

A set  $\mathcal{D} \subseteq \begin{bmatrix} V \\ k \end{bmatrix}$  is called a **simple  $t$ -( $n, k, \lambda$ ) $_q$  design**, if every  $T \in \begin{bmatrix} V \\ t \end{bmatrix}$  is contained in exactly  $\lambda$  elements of  $\mathcal{D}$ .

- ▶ set case  $q = 1$ : **combinatorial design**
- ▶  $q$ -analog case  $q \geq 2$ : **subspace design**

## Reformulation in characteristic functions

- ▶ Let  $\mathbf{x}_T$  be characteristic function of **pencil**  $\{K \in \begin{bmatrix} V \\ k \end{bmatrix} \mid T \subseteq K\}$ .
- ▶ For  $f, g : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$   
fix standard inner product  $\langle f, g \rangle = \sum_{K \in \begin{bmatrix} V \\ k \end{bmatrix}} f(K)g(K)$ .
- ▶ Note that  $\#(\mathcal{F} \cap \mathcal{G}) = \langle \chi_{\mathcal{F}}, \chi_{\mathcal{G}} \rangle$  for  $\mathcal{F}, \mathcal{G} \subseteq \begin{bmatrix} V \\ k \end{bmatrix}$ .
- ▶  $\mathcal{D}$  is simple  $t$ -( $n, k, \lambda$ ) $_q$  design  
 $\iff \langle \mathbf{x}_T, \chi_{\mathcal{D}} \rangle = \lambda$  for all  $T \in \begin{bmatrix} V \\ t \end{bmatrix}$ .
- ▶  $\rightsquigarrow$  generalization to **real designs**.

## Generalized definition: Real design

A function  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$  is called a **real  $t$ -( $n, k, \lambda$ ) $_q$  design**, if  $\langle \mathbf{x}_T, f \rangle = \lambda$  for all  $T \in \begin{bmatrix} V \\ t \end{bmatrix}$ .

- ▶  $f$  **null design** or **trade** if  $\lambda = 0$ .
- ▶  $f$  **signed design** if  $\text{im}(f) \subseteq \mathbb{Z}$ .
- ▶  $f$  **design** or **possibly non-simple design** if  $\text{im}(f) \subseteq \mathbb{N}$ .  
(Idea: simple design, but with possibly repeated blocks)
- ▶  $f$  (characteristic function of) simple design  
 $\iff \text{im}(f) \subseteq \{0, 1\} \iff f$  **Boolean**.

## Further reformulation

- ▶ Observation:  
Functions  $\mathbf{x}_T$  (interpreted as vectors)  
are the rows of incidence matrix  $W^{(tk)}$ .
- ▶ Therefore:  
 $f$  real  $t$ -( $n, k, \lambda$ ) $_q$  design  $\iff W^{(tk)}f = \lambda \mathbf{1}$ .
- ▶ In particular:  
 $f$  real  $t$ -( $n, k, 0$ ) $_q$  null design  $\iff W^{(tk)}f = \mathbf{0}$ .

## Geometric property, basic version

- ▶ For  $\lambda \in \mathbb{R}$  let  $U_\lambda :=$  set of real  $t$ -( $n, k, \lambda$ ) $_q$  design.
- ▶ Just seen:  $U_0 = \ker W^{(tk)}$ .
- ▶ Set of functions with  $A_t$  was  $V_t = \text{rowsp } W^{(tk)}$ .

$$\implies V_t = U_0^\perp$$

## What did we get?

- ▶ Established a connection to designs.
- ▶ Concept known as Delsarte's **design orthogonality**.
- ▶ Compared to prototype  
“constant intersection with all spreads”:  
Want similar property for  $\lambda \neq 0$ !

## Geometric property, version II

- ▶ Fix  $\lambda \in \mathbb{R}$ .
- ▶ Scaled complete design  $\frac{\lambda}{\lambda_{\max}} \cdot \mathbf{1}$  is real  $t$ -( $n, k, \lambda$ ) $_q$  design.
- ▶ As solution of linear equation system  $W^{(tk)}f = \lambda \mathbf{1}$ :

$$U_\lambda = \frac{\lambda}{\lambda_{\max}} \cdot \mathbf{1} + \underbrace{\ker W^{(tk)}}_{=U_0=V_t^\perp}.$$

▶  $\Rightarrow$

$$U_\lambda = \left\{ \delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#f \text{ for all } f \in V_t \right\} \quad \text{and}$$

$$V_t = \left\{ f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#f \text{ for all } \delta \in U_\lambda \right\} \quad \text{Vers. II}$$

(with  $\#f = \sum_{K \in \begin{bmatrix} V \\ k \end{bmatrix}} f(K) = \langle f, \mathbf{1} \rangle$ , motivated by  $\#\mathcal{F} = \#\chi_{\mathcal{F}}$ )

- ▶ Still room for improvement:
  - ▶ Not happy about “For all **real** ... designs”.  
 $\rightsquigarrow$  enough to look at **basis** of  $U_\lambda$ .
  - ▶ Allow mixed values of  $\lambda$ .

## Example

- ▶  $q = 1, n = 6, k = 3, t = 2 \rightsquigarrow \lambda_{\max} = \begin{bmatrix} 6-2 \\ 3-2 \end{bmatrix} = 4$ .
- ▶ Baby example:  $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$ , seen:  $\chi_{\mathcal{F}} \in V_2$ .
- ▶ Geometric property  $\implies$  For each  $2-(6, 3, 2)_1$  design:

$$\langle \chi_{\mathcal{F}}, \delta \rangle = \frac{\lambda}{\lambda_{\max}} \cdot \#\chi_{\mathcal{F}} = \frac{2}{4} \cdot 2 = 1.$$

- ▶  $\implies$  Each **simple**  $2-(6, 3, 2)_1$  design  $\mathcal{D}$  contains **exactly one** of the blocks  $\{1, 2, 3\}$  and  $\{4, 5, 6\}$ .
- ▶  $\rightsquigarrow \mathcal{D}$  is **anti-complementary**.
- ▶ Can also be shown using **intersection numbers**.

## Geometric property, toolbox version

- ▶  $U_*$  := set of all real  $t$ -( $v, k, \lambda$ ) $_q$  designs with arbitrary value  $\lambda \in \mathbb{R}$ .
- ▶ By scaled complete designs:  $U_* = U_0 + \langle \mathbf{1} \rangle_{\mathbb{R}}$ .
- ▶ Lemma (Toolbox version of geometric property).  
Let  $\Delta \subseteq U_*$ . Then

$$V_t = \left\{ f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R} \mid \langle f, \delta \rangle = \frac{\lambda_\delta}{\lambda_{\max}} \cdot \#f \text{ for all } \delta \in \Delta \right\}$$
$$\iff \langle \Delta \cup \{\mathbf{1}\} \rangle_{\mathbb{R}} = U_*$$

**Proof.** Dimension argument. Use that  $W^{(tk)}$  has full rank (Set case: Gottlieb 1966,  $q$ -analog case: Kantor 1972)

- ▶ Question: Suitable sets  $\Delta$ ?

## Lemma

Let  $\Delta$  be

- (a) the set of all signed  $t$ -( $n, k, 0$ ) $_q$  null designs or
- (b) the set of all possibly non-simple  $t$ -( $n, k, \lambda$ ) $_q$  designs

Then  $U_* = \langle \Delta \cup \{\mathbf{1}\} \rangle_{\mathbb{R}}$ .

## Proof.

Part (a).

- ▶ entries of  $W^{(tk)}$  are in  $\mathbb{Q}$ .
- ▶  $\implies U_0 = \ker W^{(tk)}$  has rational basis.
- ▶ Multiply by common denominators  $\rightsquigarrow$  integral basis  $B$ .
- ▶  $\implies B \subseteq \Delta$  and  $\langle B \cup \{\mathbf{1}\} \rangle_{\mathbb{R}} = U_*$ .

Part (b).

- ▶ Start with  $B$ .
- ▶ Add suitable integral multiples of  $\mathbf{1}$   
 $\rightsquigarrow$  non-negative integral set  $B'$ .
- ▶  $\implies B' \subseteq \Delta$  and  $\langle B' \cup \{\mathbf{1}\} \rangle_{\mathbb{R}} = U_*$ .



We arrive at:

## Theorem

Let  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$ . The following are equivalent.

(i) *Algebraic property:*  $f \in V_t$ .

*Geometric properties:*

(ii) There is a constant  $c \in \mathbb{R}$  such that  $\langle f, \delta \rangle = \lambda_\delta c$   
for all *real*  $t$ -( $n, k, *$ ) $_q$  designs  $\delta$  with  $\lambda \in \mathbb{R}$ .

(iii)  $\langle f, \delta \rangle = 0$   
for all *signed*  $t$ -( $n, k, 0$ ) $_q$  *null* designs  $\delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{Z}$ .

(iv) There is a constant  $c \in \mathbb{R}$  such that  $\langle f, \delta \rangle = \lambda_\delta c$   
for all *possibly non-simple*  $t$ -( $n, k, *$ ) $_q$  designs  $\delta : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{N}$ .

The constant in properties (ii) and (iv) necessarily equals

$$c = \frac{1}{\lambda_{\max}} \cdot \#f.$$

## Geometric property: Discussion

- ▶ Tempting: Is the following a suitable geometric property?

“There is a constant  $c \in \mathbb{R}$  such that  $\langle f, \delta \rangle = \lambda c$  for all **simple**  $t$ -( $n, k, *$ ) $_q$  designs”

- ▶ By toolbox version: If and only if  $\langle \{\text{simple } t\text{-(}n, k, *)_q \text{ designs}\} \rangle_{\mathbb{R}} = U_*$  (**richness cond**)
- ▶ Unfortunately: Not always true.

**Counterexample.**  $q = 1, n = 10, k = 5, t = 4$ .

By integrality conditions: All simple  $4$ -( $10, 5, *$ ) $_1$  are trivial.

$\implies \dim \langle \{\text{simple } 4\text{-(}10, 5, *)_1 \text{ designs}\} \rangle_{\mathbb{R}} = 1$ , too small!

- ▶ **Research problem.** (probably hard!)

Classify the parameters  $(q, n, k, t)$  where the richness condition holds.

## The Degree

► Fix  $k \in \{0, \dots, n\}$  and  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$ .

► Lemma.

$$\{\mathbf{1}\} = V_0 \subsetneq V_1 \subsetneq \dots \subsetneq V_k = V.$$

Proof.  $W^{(ij)} W^{(jk)} \sim W^{(ik)}$  for  $0 \leq i \leq j \leq k$ .

► Definition.

Degree  $\deg(f) :=$  smallest  $t$  such that  $f \in V_t$ .

## Example

► Functions  $f$  of degree 0  
are the scalar functions  $f = \lambda \mathbf{1}$  with  $\lambda \in \mathbb{R}$ .

► Baby example  $\mathcal{F} = \{\{1, 2, 3\}, \{4, 5, 6\}\}$ .

In  $V = \{1, 2, 3, 4, 5, 6\}$  we have  $\deg(\mathcal{F}) := \deg(\chi_{\mathcal{F}}) = 2$ .

► Seen:  $\chi_{\mathcal{F}} \in V_2$ .

► Exercise:  $\chi_{\mathcal{F}} \notin V_1$ .

► In  $V = \{1, 2, 3, 4, 5, 6, 7\}$  we have  $\deg(\mathcal{F}) = 3$ .

⇒ Ambient space  $V$  matters!

## The Degree (cont.)

► **Remember.** Rows of  $W^{(tk)}$  are the  $t$ -pencils  $\mathbf{x}_T$ .

►  $\rightsquigarrow$  **Alternative characterization of degree.**

$\deg(f)$  is smallest  $t$

such that  $f$  is a linear combination of  $t$ -pencils  $\mathbf{x}_T$ .

The (unique) coefficients are called  **$t$ -weights**  $\text{wt}_f(T)$  of  $f$ :

$$f = \sum_{T \in \binom{V}{t}} \text{wt}_f(T) \mathbf{x}_T$$

## Lemma

(a)  $\deg(\lambda f) \leq \deg(f)$  *with equality iff*  $\lambda \neq 0$ .

(b)  $\deg(f + g) \leq \max(\deg(f), \deg(g))$ .

(c)  $\deg(fg) \leq \deg(f) + \deg(g)$ .

## Proof.

Parts (a), (b): easy. Part (c): Use weights &  $\deg \mathbf{x}_T \leq \text{rk } T$ . □

## Dualization

- ▶ Fix anti-isomorphism  $\perp$  of the lattice  $\mathcal{L}(V)$ .
  - ▶ Set case: Set complement.
  - ▶  $q$ -analog case: Perp wrt non-degenerate bilinear form.
- ▶ Induces **dual map** of  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{R}$ :

$$\textcolor{red}{f}^\perp : \begin{bmatrix} V \\ n-k \end{bmatrix} \rightarrow \mathbb{R}, \quad U \mapsto f(U^\perp)$$

- ▶ Effect of dualization on the degree?

## Theorem

(a)  $\deg f^\perp = \deg f$ .

(b) For  $i \in \{0, \dots, \deg f\}$ , the *i-weight distribution* of  $f^\perp$  is

$$\text{wt}_{f^\perp}^{(i)}(J) = \sum_{I \in \binom{V}{i}} \gamma(n-k, i, \text{rk}(I^\perp \cap J)) \text{wt}_f^{(i)}(I)$$

where

$$\gamma(k, i, z) := \begin{cases} \delta_{z,k} & \text{if } i = k, \\ (-1)^{i-z} \frac{1}{q^{(k-i)(i-z) + \binom{i-z}{2}}} \frac{\begin{bmatrix} k-i \\ 1 \end{bmatrix}}{\begin{bmatrix} k-z \\ 1 \end{bmatrix}} \frac{1}{\begin{bmatrix} k \\ z \end{bmatrix}} & \text{otherw.} \end{cases}$$

## Proof.

- ▶ Enough to look at *pencils*  $f = \mathbf{x}_J$ .
- ▶ Set up linear equation system for the weights of  $f^\perp$ , assuming that  $\text{wt}(I)$  only depends on  $\text{rk}(I \cap J)$ .
- ▶ Equation system matrix is *triangular* with non-zero diagonal  $\implies$  invertible  $\implies$  Part (a).
- ▶ Apply negation formula &  $q$ -Vandermonde formula for Gaussian coefficients  $\rightsquigarrow$  compute solution  $\rightsquigarrow$  Part (b).

## Change of ambient space

Two elementary ways to shrink the ambient space  $V$ .

- ▶  $V \rightarrow H$  ( $H \in \begin{bmatrix} V \\ n-1 \end{bmatrix}$  hyperplane)
- ▶  $V \rightarrow V/P$  ( $P \in \begin{bmatrix} V \\ 1 \end{bmatrix}$  point)

Implication on the degree?

We start with  $V \rightarrow V/P$ .

## Theorem

Let  $1 \leq k \leq n$  and  $P \in \begin{bmatrix} V \\ 1 \end{bmatrix}$ . Then

$$\Phi : \mathbb{R}^{\begin{bmatrix} V/P \\ k-1 \end{bmatrix}} \rightarrow \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}}, \quad \Phi(f) : K \mapsto \begin{cases} f(K/P) & \text{if } P \subseteq K, \\ 0 & \text{if } P \not\subseteq K \end{cases}$$

is an injective  $\mathbb{R}$ -linear map with

$$\text{im}(\Phi) = \{g \in \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}} \mid \text{supp } g \subseteq \begin{bmatrix} V \\ k \end{bmatrix}|_P\} \quad \text{and}$$

$$\deg_V \Phi(f) = \begin{cases} 0 & \text{if } f = 0, \\ \min(\overbrace{\deg_{V/P}(f) + 1}^{\text{main case}}, n - k) & \text{otherwise.} \end{cases}$$

## Proof.

- ▶ Straightforward, except “ $\deg_V \Phi(f) \geq \deg_{V/P}(f) + 1$ ”.
- ▶ **Lemma.** In main case

For all  $g \in \text{im } \Phi$ :  $P \not\subseteq T \implies \text{wt}_g(T) = 0$ .

**Proof.** Incidence matrices of certain attenuated geometries are of full rank. (Guo, Li, Wang, 2014.)



## Theorem

Let  $1 \leq n - k \leq n$  and  $H \in \begin{bmatrix} V \\ n-1 \end{bmatrix}$ . Then

$$\psi : \mathbb{R}^{\begin{bmatrix} H \\ k \end{bmatrix}} \rightarrow \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}}, \quad \psi(f) : K \mapsto \begin{cases} f(K) & \text{if } K \subseteq H, \\ 0 & \text{if } K \not\subseteq H \end{cases}$$

is an injective  $\mathbb{R}$ -linear map with

$$\begin{aligned} \text{im}(\psi) &= \{g \in \mathbb{R}^{\begin{bmatrix} V \\ k \end{bmatrix}} \mid \text{supp } g \subseteq \begin{bmatrix} H \\ k \end{bmatrix}\} \quad \text{and} \\ \deg_V \psi(f) &= \begin{cases} 0 & \text{if } f = 0, \\ \min(\deg_H(f) + 1, k) & \text{otherwise.} \end{cases} \end{aligned}$$

## Proof.

Follows from the previous theorem by **dualization**. □

## Example (Basic sets)

- ▶ Start with “complete set”  $\begin{bmatrix} W \\ \ell \end{bmatrix}$  of degree 0.



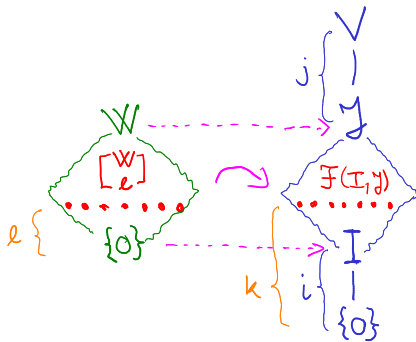
- ▶  $i$ -fold application of  $\Phi$  and  $j$ -fold application of  $\Psi$

$$\rightsquigarrow \text{basic set } \mathcal{F}(I, J) = \left\{ K \in \begin{bmatrix} V \\ k \end{bmatrix} \mid I \subseteq K \subseteq J \right\}.$$

- ▶ By theorems:  $\deg \mathcal{F}(I, J) = i + j$ .

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## Example (Basic sets (cont.))

- ▶ Basic sets  $\mathcal{F}(I, J)$  include pencils ( $j = 0$ ) and dual pencils ( $i = 0$ ).  
In particular  $\deg \mathbf{x}_I = \text{rk } I$ .
- ▶ Geometric property of  $\mathcal{F}(I, J)$



design property of  $i$ -fold derived and  $j$ -fold residual design.

## Sets and Boolean functions

- ▶ Of particular interest: **Sets**  $\mathcal{F} \subseteq \binom{V}{k}$  of low degree.
- ▶ Via characteristic functions:  
Sets correspond to **Boolean** functions  $\binom{V}{k} \rightarrow \{0, 1\}$ .

## Boolean degree 1 functions

- ▶ **Set** case:

Filmus, Ihringer 2019:

Only **basic** functions.

$\implies$  only pencils and dual pencils (since  $t = 1$ ).

- ▶  **$q$ -analog** case:

Boolean degree 1 function = **Cameron-Liebler** set of  $(k - 1)$ -spaces in  $\text{PG}(n - 1, q)$ .

**Non-basic** examples do exist.

Classification: Hard research problem.

# Computer classification

## Goal.

For  $q = 1$  and small  $n, k$ , **classify** all sets  $\mathcal{F}$  of degree  $t = 2$ .

## Strategy.

- ▶ Use “basic” **geometric property**:

$$\deg \chi_{\mathcal{F}} \leq t \iff \chi_{\mathcal{F}} \in \ker W^{(tk)}.$$

$\rightsquigarrow$  Want to find all  $\{0, 1\}$ -vectors in  $\ker W^{(tk)}$ .

- ▶ Find **integral basis** of  $\ker W^{(tk)}$ .
  - ▶ either: computationally
  - ▶ or: Use literature like  
Khosrovshahi, Ajoodani-Namini (1990):  
*A new basis for trades*
- ▶  $\rightsquigarrow$  system of linear Diophantine equations.
- ▶ Solve using **SOLVEDIOPHANT** (A. Wassermann)
- ▶ Filter out isomorphic copies.  
(action of symmetric group  $\mathfrak{S}_n$ )

# Results

$n$	$k$	size distribution	$\Sigma$
6	3	2 4 <sup>3</sup> 6 <sup>5</sup> 8 <sup>8</sup> 10 <sup>10</sup> 12 <sup>8</sup> 14 <sup>5</sup> 16 <sup>3</sup> 18	44
7	3	5 <sup>2</sup> 10 <sup>6</sup> 15 <sup>12</sup> 20 <sup>12</sup> 25 <sup>6</sup> 30 <sup>2</sup>	40
8	3	6 8 11 12 14 15 <sup>2</sup> 16 17 18 20 <sup>2</sup> 21 <sup>2</sup> 22 23 24 <sup>3</sup> 25 26 <sup>4</sup> 27 28 <sup>2</sup> ...	52
9	3	7 14 <sup>2</sup> 21 <sup>5</sup> 28 <sup>5</sup> 35 <sup>5</sup> 42 <sup>11</sup> 49 <sup>5</sup> 56 <sup>5</sup> 63 <sup>5</sup> 70 <sup>2</sup> 77	47
10	3	8 16 20 24 <sup>2</sup> 28 <sup>3</sup> 32 <sup>2</sup> 36 <sup>4</sup> 40 <sup>2</sup> 44 <sup>2</sup> 48 <sup>2</sup> 52 <sup>2</sup> 56 <sup>5</sup> 60 <sup>5</sup> 64 <sup>5</sup> 68 <sup>2</sup> 72 <sup>2</sup> ...	59
8	4	10 15 <sup>2</sup> 20 <sup>3</sup> 30 <sup>6</sup> 35 <sup>4</sup> 40 <sup>6</sup> 50 <sup>3</sup> 55 <sup>2</sup> 60	28
9	4	21 <sup>2</sup> 35 <sup>3</sup> 56 <sup>5</sup> 70 <sup>5</sup> 91 <sup>3</sup> 105 <sup>2</sup>	20
10	4	28 42 56 <sup>2</sup> 70 84 <sup>2</sup> 98 <sup>3</sup> 112 <sup>3</sup> 126 <sup>2</sup> 140 154 <sup>2</sup> 168 182	20
11	4	36 78 84 <sup>2</sup> 120 <sup>2</sup> 126 162 <sup>3</sup> 168 <sup>3</sup> 204 210 <sup>2</sup> 246 <sup>2</sup> 252 294	20
12	4	45 120 <sup>2</sup> 135 165 <sup>2</sup> 210 240 <sup>3</sup> 255 <sup>3</sup> 285 330 <sup>2</sup> 360 375 <sup>2</sup> 450	20

blue = sizes of basic sets

**Goal.** Explain divisibility pattern of the sizes!

## Theorem (Divisibility theorem)

Let  $f : \begin{bmatrix} V \\ k \end{bmatrix} \rightarrow \mathbb{Z}$  be a function of degree  $t$ . Then

$$\underbrace{\gcd \left( \begin{bmatrix} n-0 \\ k-0 \end{bmatrix}, \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}, \dots, \begin{bmatrix} n-t \\ k-t \end{bmatrix} \right)}_{=: a} \mid \#f.$$

### Proof.

► Algebraic property  $\Rightarrow \exists \mathbf{x} : \begin{bmatrix} V \\ t \end{bmatrix} \rightarrow \mathbb{R}$  with  $\mathbf{x}^\top W^{(tk)} = f^\top$ . (1)

► Complete design:  $W^{(tk)} \cdot \mathbf{1} = \lambda_{\max} \cdot \mathbf{1}$  (2)

► Design theory:

parameters  $t$ - $(n, k, \lambda_{\min})_q$  with  $\lambda_{\min} = \frac{\lambda_{\max}}{a}$  are admissible.

►  $\Rightarrow \exists$  signed  $t$ - $(n, k, \lambda_{\min})_q$  design  $\delta \Rightarrow W^{(tk)}\delta = \lambda_{\min} \cdot \mathbf{1}$  (3)

► Set case: Wilson, "The necessary conditions for  $t$ -designs are sufficient for something" (1973).

►  $q$ -analog case: Ray-Chaudhuri, Singhi (1989).

► Left multiplication of (2) and (3) by  $\mathbf{x}^\top$ , using (1)

$$\implies \#f = \lambda_{\max} \cdot \#\mathbf{x} \quad \text{and} \quad \langle f, \delta \rangle = \lambda_{\min} \cdot \#\mathbf{x}$$

►  $\implies \#f = \underbrace{a \cdot \langle f, \delta \rangle}_{\in \mathbb{Z}} \in \mathbb{Z}.$



# Compare with the results

$$q = 1, t = 2 \implies a = \gcd\left(\binom{n}{k}, \binom{n-1}{k-1}, \binom{n-2}{k-2}\right).$$

$n$	$k$	size distribution	$a$
6	3	$2 \cdot 4^3 \cdot 6^5 \cdot 8^8 \cdot 10^{10} \cdot 12^8 \cdot 14^5 \cdot 16^3 \cdot 18$	2
7	3	$5^2 \cdot 10^6 \cdot 15^{12} \cdot 20^{12} \cdot 25^6 \cdot 30^2$	5
8	3	$6 \cdot 8 \cdot 11 \cdot 12 \cdot 14 \cdot 15^2 \cdot 16 \cdot 17 \cdot 18 \cdot 20^2 \cdot 21^2 \cdot 22 \cdot 23 \cdot 24^3 \cdot 25 \cdot 26^4 \cdot 27 \cdot 28^2 \dots$	1
9	3	$7 \cdot 14^2 \cdot 21^5 \cdot 28^5 \cdot 35^5 \cdot 42^{11} \cdot 49^5 \cdot 56^5 \cdot 63^5 \cdot 70^2 \cdot 77$	7
10	3	$8 \cdot 16 \cdot 20 \cdot 24^2 \cdot 28^3 \cdot 32^2 \cdot 36^4 \cdot 40^2 \cdot 44^2 \cdot 48^2 \cdot 52^2 \cdot 56^5 \cdot 60^5 \cdot 64^5 \cdot 68^2 \cdot 72^2 \dots$	4
8	4	$10 \cdot 15^2 \cdot 20^3 \cdot 30^6 \cdot 35^4 \cdot 40^6 \cdot 50^3 \cdot 55^2 \cdot 60$	5
9	4	$21^2 \cdot 35^3 \cdot 56^5 \cdot 70^5 \cdot 91^3 \cdot 105^2$	7
10	4	$28 \cdot 42 \cdot 56^2 \cdot 70 \cdot 84^2 \cdot 98^3 \cdot 112^3 \cdot 126^2 \cdot 140 \cdot 154^2 \cdot 168 \cdot 182$	14
11	4	$36 \cdot 78 \cdot 84^2 \cdot 120^2 \cdot 126 \cdot 162^3 \cdot 168^3 \cdot 204 \cdot 210^2 \cdot 246^2 \cdot 252 \cdot 294$	6
12	4	$45 \cdot 120^2 \cdot 135 \cdot 165^2 \cdot 210 \cdot 240^3 \cdot 255^3 \cdot 285 \cdot 330^2 \cdot 360 \cdot 375^2 \cdot 450$	15

Perfect fit!

## Parameter of Cameron-Liebler sets of $k$ -spaces

- ▶ Consider  $q$ -analog case  $q \geq 2$ .
- ▶ For sets  $\mathcal{F}$  of degree  $t = 1$  define  
**parameter  $x$**   $:= \#\mathcal{F} / \begin{bmatrix} n-1 \\ k-1 \end{bmatrix} \in \mathbb{Q}$
- ▶ **Corollary** of divisibility theorem.

$$\frac{q^k - 1}{q^{\gcd(n,k)} - 1} \cdot x \in \mathbb{Z},$$

restricting denominator of fraction  $x$  in canceled form.

### Example

- ▶  $k \mid n \implies x \in \mathbb{Z}$ .  
Already known: Blokhuis, De Boeck, D'haeseleer (2019).
- ▶  $n$  and  $k$  coprime  $\implies (1 + q + \dots + q^{k-1}) \cdot x \in \mathbb{Z}$ .
- ▶  $k = 4, n \equiv 2 \pmod{4} \implies (1 + q^2) \cdot x \in \mathbb{Z}$ .

## The paired construction

- ▶ Construction for the **set case**  $q = 1$  only.
- ▶ **Idea.** Disjoint union of two “opposite” basic sets.
- ▶ Let  $I, J \subseteq V$  be disjoint, not both empty. Define

$$\mathcal{P}(I, J) := \mathcal{F}(I, J^c) \uplus \mathcal{F}(J, I^c)$$

- ▶ Clear:

$$\deg \mathcal{P}(I, J) \leq \min(\#I + \#J, k) \quad (\text{trivial bound}).$$

- ▶ Will see: There are cases with a **strict** “ $<$ ”!

## Example

- $V = \{1, \dots, 6\}$ ,  $k = 3$ ,  $I = \emptyset$ ,  $J = \{4, 5, 6\}$ ,

$$\begin{aligned} & \mathcal{P}(\emptyset, \{4, 5, 6\}) \\ &= \mathcal{F}(\emptyset, \{1, 2, 3\}) \uplus \mathcal{F}(\{4, 5, 6\}, \{1, 2, 3, 4, 5, 6\}) \\ &= \{\{1, 2, 3\}, \{4, 5, 6\}\} \quad \text{Baby example!} \end{aligned}$$

- Already seen:

$\deg \mathcal{P}(\emptyset, \{1, 2, 3\}) = 2$ , beating the trivial bound “ $\leq 3$ ”!

## Example

- $V = \{1, \dots, 7\}$ ,  $k = 3$ ,  $I = \{1\}$ ,  $J = \{6, 7\}$ .

$$\begin{aligned} & \mathcal{P}(\{1\}, \{6, 7\}) \\ &= \mathcal{F}(\{1\}, \{1, 2, 3, 4, 5\}) \uplus \mathcal{F}(\{6, 7\}, \{2, 3, 4, 5, 6, 7\}) \\ &= \{\{1, 2, 3\}, \{1, 2, 4\}, \{1, 2, 5\}, \{1, 3, 4\}, \{1, 3, 5\}, \{1, 4, 5\}, \\ & \quad \{2, 6, 7\}, \{3, 6, 7\}, \{4, 6, 7\}, \{5, 6, 7\}\} \end{aligned}$$

- $\deg(\mathcal{P}(\{1\}, \{6, 7\})) = 2$ , beating the trivial bound.

## Theorem

Let  $q = 1$ ,  $I, J \subseteq V$  disjoint,  $i = \#I$ ,  $j = \#J$ ,  $k \leq \frac{n}{2}$ ,  
 $i \leq k \leq n - i$ ,  $j \leq k \leq n - j$ .

In the cases

- (a)  $i + j \leq k$  and  $i + j$  odd;
- (b)  $i + j \geq k$  and  $k$  odd and  $n = 2k$

we have

$$\deg \mathcal{P}(I, J) \leq \min(i + j, k) - 1.$$

## Proof (Idea).

Part (a): Write  $\chi_{\mathcal{P}(I, J)}$  as an integer linear combination of basic functions of degree  $i + j - 1$ .

Part (b):

- ▶ Use  $\mathcal{P}(X, Y) = \mathcal{P}(X \uplus \{x\}, Y) \uplus \mathcal{P}(X, Y \uplus \{x\})$   
(where  $X, Y, \{x\}$  are pairwise disjoint)
- ▶ Moving elements from  $J$  to  $I \rightsquigarrow \deg \mathcal{P}(I, J) \leq \deg \mathcal{P}(K, J')$
- ▶  $\mathcal{P}(K, J') = \mathcal{P}(K, \emptyset) \implies$  Back in Case (a).

## Theorem

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## Conjecture

Statement of Theorem is **best possible**.

- In fact always equality

$$\deg \mathcal{P}(I, J) = \min(i + j, k) - 1.$$

- In all cases not covered by (a) and (b),  
the trivial bound is sharp:

$$\deg \mathcal{P}(I, J) = \min(i + j, k).$$

## Small sets of degree $t$

- ▶ Natural question.

Smallest size  $m_q(n, k, t)$  of a non-empty set of degree  $\leq t$ ?

- ▶ From  $\deg \mathbf{x}_T = t$  we get

$$m_q(n, k, t) \leq \begin{bmatrix} n - t \\ k - t \end{bmatrix}. \quad (*)$$

- ▶ Bound  $(*)$  is always sharp for  $t = 1$ .

- ▶ Set case: Filmus, Ihringer (2019).

- ▶  $q$ -analog case: Blokhuis, De Boeck, D'haeseleer (2019).

- ▶ For  $q = 1$ ,  $n = 2k$ ,  $t \geq 2$  even,  $i = 0$  and  $j = t + 1$ , the paired construction beats bound  $(*)$ !

## Corollary

Let  $t \in \{0, \dots, k - 1\}$  be even. Then

$$m_1(2k, k, t) \leq 2 \cdot \binom{2k - t - 1}{k}.$$

## Open problems

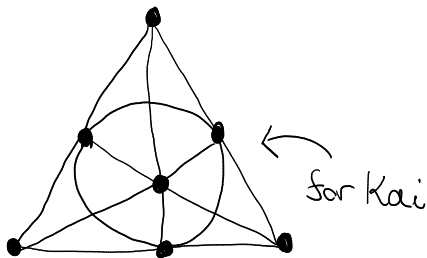
- ▶ Many!
- ▶ For fixed  $(q, n, k, t)$ , characterize the **sizes** of degree  $t$  sets.
  - ▶ Smallest,
  - ▶ second smallest,
  - ▶ gaps,
  - ▶ etc.
- ▶ Further investigate and exploit **relationship**  
degree  $t$  functions  $\longleftrightarrow$   $t$ -designs.

Which results can be translated?

- ▶ Maybe most important:  
Better **name** for the studied objects.
  - ▶ “dual designs”?  $\longrightarrow$  ambiguous.
  - ▶ Something involving “Cameron-Liebler”?
  - ▶ other ideas?



# Thank you!



Slides will be uploaded at

<https://mathe2.uni-bayreuth.de/michaelk/>