



# Rational Points on Curves of Genus 2

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## Curves of Genus 2

A **curve of genus 2** over  $\mathbb{Q}$  is given by an equation

$$C : y^2 = f_6x^6 + f_5x^5 + f_4x^4 + f_3x^3 + f_2x^2 + f_1x + f_0$$

with  $f_j \in \mathbb{Z}$ , such that  $(f_6, f_5) \neq (0, 0)$

and the polynomial on the right does not have multiple roots.

A **rational point** on this curve  $C$

is a pair of rational numbers  $(\xi, \eta)$  satisfying the equation.

In addition, there can be rational points **“at infinity”**, corresponding to the square roots of  $f_6$  in  $\mathbb{Z}$ .

We denote the set of rational points on  $C$  by  $C(\mathbb{Q})$ .

**Theorem** (Mordell’s Conjecture, proved by Faltings).  $C(\mathbb{Q})$  is **finite**.

# The Questions

Consider curves of **genus 2** over  $\mathbb{Q}$ :

$$C : y^2 = f_6x^6 + f_5x^5 + f_4x^4 + f_3x^3 + f_2x^2 + f_1x + f_0$$

with  $f_j \in \mathbb{Z}$ ,  $|f_j| \leq N$ .

## Question.

What can we say about  $C(\mathbb{Q})$ , the set of **rational points** on  $C$ , as  $N$  grows?

- How do we determine  $C(\mathbb{Q})$ ?
- How large is  $C(\mathbb{Q})$  on average?
- How large can the points get?
- How many curves have many rational points?
- How are the sizes of the points distributed?

## Heuristics (1)

The condition that the point  $(\frac{a}{b}, \frac{c}{b^3})$  is on  $C$  translates into a **linear condition** on the coefficients  $f_j$ :

$$a^6 f_6 + a^5 b f_5 + a^4 b^2 f_4 + a^3 b^3 f_3 + a^2 b^4 f_2 + a b^5 f_1 + b^6 f_0 = c^2$$

The curves satisfying this correspond to points in the intersection of a coset of a 6-dimensional **lattice** in  $\mathbb{R}^7$  with a **cube** of side length  $2N$ .

We can **estimate** the size of this set by the **volume** of the corresponding slice of the cube, divided by the **covolume** of the lattice.

We obtain for the average number of points with  $x = \frac{a}{b}$ :

$$\mathbb{E}_N(\#C(\mathbb{Q})_x) \sim \frac{\gamma(x)}{\sqrt{N}} \quad \text{as } N \rightarrow \infty.$$

with  $\gamma(x)$  of order  $H(x)^{-3}$ , where  $H(x) = \max\{|a|, |b|\}$  is the **height** of  $x$ .

## Heuristics (2)

We let

$$\gamma(H) = \sum_{x \in \mathbb{P}^1(\mathbb{Q}), H(x) \leq H} \gamma(x) = \gamma - O\left(\frac{1}{H}\right)$$

where  $\gamma = \lim_{H \rightarrow \infty} \gamma(H) < \infty$ .

If  $C(\mathbb{Q})_H = \{(x, y) \in C(\mathbb{Q}) : H(x) \leq H\}$ , this gives

$$\mathbb{E}_N(\#C(\mathbb{Q})_H) \sim \frac{\gamma(H)}{\sqrt{N}} \quad \text{as } N \rightarrow \infty.$$

**Corollary.**

$$\liminf_{N \rightarrow \infty} \sqrt{N} \cdot \mathbb{E}_N(\#C(\mathbb{Q})) \geq \gamma.$$

**Conjecture.**

$$\lim_{N \rightarrow \infty} \sqrt{N} \cdot \mathbb{E}_N(\#C(\mathbb{Q})) = \gamma.$$

## Heuristics (3)

Let  $\pi : C \rightarrow \mathbb{P}^1$  be the  $x$ -coordinate map.

For  $P \in C(\mathbb{Q})$ , we write  $H(P) = H(\pi(P))$ .

The linear conditions for up to seven  $x$ -coordinates are **independent**, so we expect (at least for  $k \leq 7$ )

$$\text{Prob}_N(\#\pi(C(\mathbb{Q})) \geq k) \sim \frac{c_k}{N^{k/2}}.$$

We also expect that (using  $\gamma - \gamma(H) \approx \beta/H$ )

$$\#\{C : \exists P \in C(\mathbb{Q}), H(P) \geq H\} \approx \beta \frac{N^{13/2}}{H},$$

which leads to the

**Conjecture.**  $\max\{H(P) : P \in \cup_C C(\mathbb{Q})\} \ll N^{13/2+\varepsilon}$ .

(Note: This is **polynomial** in  $N$ , in contrast to curves of genus 1.)

# How To Find the Points

In order to test these heuristics, we need to **find**  $C(\mathbb{Q})$  for all curves  $C$  (with  $|f_j| \leq N$ ).

We can **search** for points on  $C$ ; this is feasible for heights up to  $10^4$  or maybe  $10^5$  (using `ratpoints`, or `Points()` in MAGMA; complexity of order  $H^2$ ).

If we do not find any points, we can try to **prove** that  $C(\mathbb{Q}) = \emptyset$ :

- **local obstruction**:  $C(\mathbb{R}) = \emptyset$  or  $C(\mathbb{Q}_p) = \emptyset$  for some  $p$
- **2-cover descent**:  $\text{Sel}^{(2)}(C/\mathbb{Q}) = \emptyset$
- **Mordell-Weil Sieve** (see below)

## Example.

For  $N = 3$ , there are **196 171** isomorphism classes of curves, of which **137 490** have rational points, and **58 681** don't. (Bruin-Stoll)

# The Jacobian Variety

The points on an elliptic curve form a **group** in a natural way. This helps tremendously when studying such curves.

The points on a curve of genus 2 **do not** form a group in a natural way.

However, we can **embed**  $C$  into a 2-dimensional variety  $J$  whose points **do** form a group in a natural way.

This variety  $J$  is called the **Jacobian variety** of  $C$ .

Its rational points form a group  $J(\mathbb{Q})$ .

Weil generalised a theorem of Mordell's on elliptic curves and showed that  $J(\mathbb{Q})$  is a **finitely generated** abelian group.

We call  $J(\mathbb{Q})$  the **Mordell-Weil group** of  $J$ .



# Determining the Mordell-Weil Group

We can use the Mordell-Weil group to get information on  $C(\mathbb{Q})$ .

For this, we need to **compute generators** of  $J(\mathbb{Q})$ ;

in particular, we need to find its **rank**  $\text{rank } J(\mathbb{Q}) = \dim_{\mathbb{Q}} J(\mathbb{Q}) \otimes_{\mathbb{Z}} \mathbb{Q}$ .

- **2-descent** on  $J$  gives **upper bound** for the rank
- **Search** for points on  $J$  gives **lower bound** for the rank
- Use canonical **height**  $\hat{h}$  to saturate the known subgroup

(Algorithms are implemented in MAGMA.)

## Potential Problems.

- Upper bound may not be tight
- Some generators may be too large to be found

# Using the Mordell-Weil Group

Let  $\iota : C \rightarrow J$  be an embedding.

Then  $\iota(C(\mathbb{Q})) = J(\mathbb{Q}) \cap \iota(C)$ .

Let  $r = \text{rank } J(\mathbb{Q})$ .

## Method 1.

We have  $\hat{h}(\iota(P)) \leq \log H(P) + c$ .

**Enumerate** all points  $Q \in J(\mathbb{Q})$  with  $\hat{h}(Q) \leq \log H + c$

(working with the Mordell-Weil **lattice**  $(J(\mathbb{Q})/\text{torsion}, \hat{h})$ )

and check for each  $Q$  if  $Q \in \iota(C)$ .

- Complexity of order  $\frac{\#J(\mathbb{Q})_{\text{tors}}}{\sqrt{\text{Reg}_J}} \frac{\pi^{r/2} (\log H + c)^{r/2}}{(r/2)!} = O((\log H)^{r/2})$
- Check  $Q \in \iota(C)$  first **modulo**  $p$  for many  $p$ .
- Can do  $H = 10^{100}$  for  $r = 3, 4, 5$ .

# The Mordell-Weil Sieve (1)

## Method 2.

Let  $S$  be a **finite set of primes of good reduction** for  $C$ .  
Consider the following diagram.

$$\begin{array}{ccccc}
 C(\mathbb{Q}) & \xrightarrow{\iota} & J(\mathbb{Q}) & \longrightarrow & J(\mathbb{Q})/nJ(\mathbb{Q}) \\
 \downarrow & & \downarrow & & \downarrow \beta \\
 \prod_{p \in S} C(\mathbb{F}_p) & \xrightarrow{\iota} & \prod_{p \in S} J(\mathbb{F}_p) & \longrightarrow & \prod_{p \in S} J(\mathbb{F}_p)/nJ(\mathbb{F}_p) \\
 & & \searrow \alpha & & \\
 & & & & 
 \end{array}$$

We **can compute** the maps  $\alpha$  and  $\beta$ .

If the image of the known rational points on  $C$  coincides with  $\beta^{-1}(\text{im}(\alpha))$ , then any **unknown point** must have  **$\log H(P) \gg n^2$** .

(If  $\beta^{-1}(\text{im}(\alpha)) = \emptyset$ , i.e.,  **$\text{im}(\alpha) \cap \text{im}(\beta) = \emptyset$** , this proves that  **$C(\mathbb{Q}) = \emptyset$** .)

## The Mordell-Weil Sieve (2)

A carefully optimized version of the Mordell-Weil sieve works very well when  $r = 2$  and allows us to reach  $H = 10^{1000}$  (or more) in this case.

### Example (Bruin-Stoll).

For the 1492 curves  $C$  for  $N = 3$  without rational points that do not have a local obstruction or a 2-cover obstruction, a Mordell-Weil sieve computation proves that  $C(\mathbb{Q}) = \emptyset$ .

(For 42 curves,

we need to assume the Birch and Swinnerton-Dyer Conjecture.)

## A Refinement

Taking  $n$  as a **multiple of  $N$** ,  
the Mordell-Weil sieve gives us a way of proving  
that a given **coset** of  $NJ(\mathbb{Q})$  does not meet  $\iota(C)$ .

### **Conjecture.**

If  $(Q + NJ(\mathbb{Q})) \cap \iota(C) = \emptyset$ , then there are  $n \in N\mathbb{Z}$  and  $S$  such that  
the Mordell-Weil sieve with these parameters **proves** this fact.

So if we can find an  $N$  that **separates** the rational points on  $C$ ,  
i.e., such that the composition  $C(\mathbb{Q}) \xrightarrow{\iota} J(\mathbb{Q}) \rightarrow J(\mathbb{Q})/NJ(\mathbb{Q})$  is **injective**,  
then we **can effectively determine  $C(\mathbb{Q})$**  if the Conjecture holds for  $C$ :

For each coset of  $NJ(\mathbb{Q})$ , we either **find** a point on  $C$  mapping into it,  
or we **prove** that there is no such point.

# Chabauty's Method

Chabauty's method allows us to **compute** a separating  $N$  when  $r = 1$ .

Let  $p$  be a prime of good reduction for  $C$ . There is a pairing

$$\Omega_J^1(\mathbb{Q}_p) \times J(\mathbb{Q}_p) \longrightarrow \mathbb{Q}_p, \quad (\omega, R) \longmapsto \int_0^R \omega.$$

Since  $\text{rank } J(\mathbb{Q}) = 1 < 2 = \dim_{\mathbb{Q}_p} \Omega_J^1(\mathbb{Q}_p)$ , there is a differential  $0 \neq \omega_p \in \Omega_C(\mathbb{Q}_p) \cong \Omega_J^1(\mathbb{Q}_p)$  that **kills**  $J(\mathbb{Q}) \subset J(\mathbb{Q}_p)$ .

## Theorem.

If the reduction  $\bar{\omega}_p$  **does not vanish on**  $C(\mathbb{F}_p)$  and  $p > 2$ , then each residue class mod  $p$  contains **at most one** rational point.

This implies that  $N = \#J(\mathbb{F}_p)$  is **separating**.

# Chabauty + MW Sieve

We can easily compute  $\bar{\omega}_p$ .

Heuristically (at least if  $J$  is **simple**),  
we expect to find **many**  $p$  satisfying the condition.

In practice, such  $p$  are easily found;  
the Mordell-Weil sieve computation then **determines**  $C(\mathbb{Q})$  very quickly.

## **Note.**

If  $r = 0$ , then  $C(\mathbb{Q}) = \iota^{-1}(J(\mathbb{Q})_{\text{tors}})$  is easily computed.

# Summary: Finding Points

For the 137 490 curves with rational points ( $N = 3$ ):

- $r = 0$  (14 010 curves):  $C(\mathbb{Q})$  is determined.
- $r = 1$  (46 575 curves):  $C(\mathbb{Q})$  is determined.
- $r = 2$  (52 227 curves):  $C(\mathbb{Q})_H$  is determined for  $H = 10^{1000}$ .
- $r = 3$  (22 343 curves):  $C(\mathbb{Q})_H$  is determined for  $H = 10^{100}$ .
- $r = 4$  ( 2 318 curves):  $C(\mathbb{Q})_H$  is determined for  $H = 10^{100}$ .
- $r = 5$  ( 17 curves):  $C(\mathbb{Q})_H$  is determined for  $H = 10^{100}$ .

(Caveat: In some cases when  $r$  is less than the Selmer rank, the rank is not yet proved to be correct.)

In addition, all points up to  $H = 2^{14} - 1 = 16\,383$  were computed on all curves with  $N \leq 10$ . (This was done using `ratpoints`.)



## Some Records (1)

Maximal **Number** of Points ( $H(P) \leq 16383$  for  $N \geq 4$ )

size of curves $N$	1	2	3	4	5	6	7–10
max. $\#C(\mathbb{Q})$	18	24	26	36	38	44	52

**Example.**  $y^2 = x^6 - 2x^5 + x^4 - 5x^3 - 5x^2 + 7x + 4$

has **52 points**, with  $x$ -coordinates

$$\infty, -3, -1, 0, 1, 3, 4, -\frac{5}{2}, -\frac{1}{2}, \frac{1}{2}, \frac{5}{2}, \frac{11}{2}, -\frac{4}{3}, -\frac{1}{3}, -\frac{9}{4}, -\frac{1}{4},$$

$$\frac{13}{5}, \frac{5}{6}, -\frac{3}{7}, -\frac{63}{11}, \frac{96}{11}, -\frac{7}{15}, \frac{19}{20}, -\frac{265}{111}, -\frac{7}{369}, \frac{199}{504}.$$

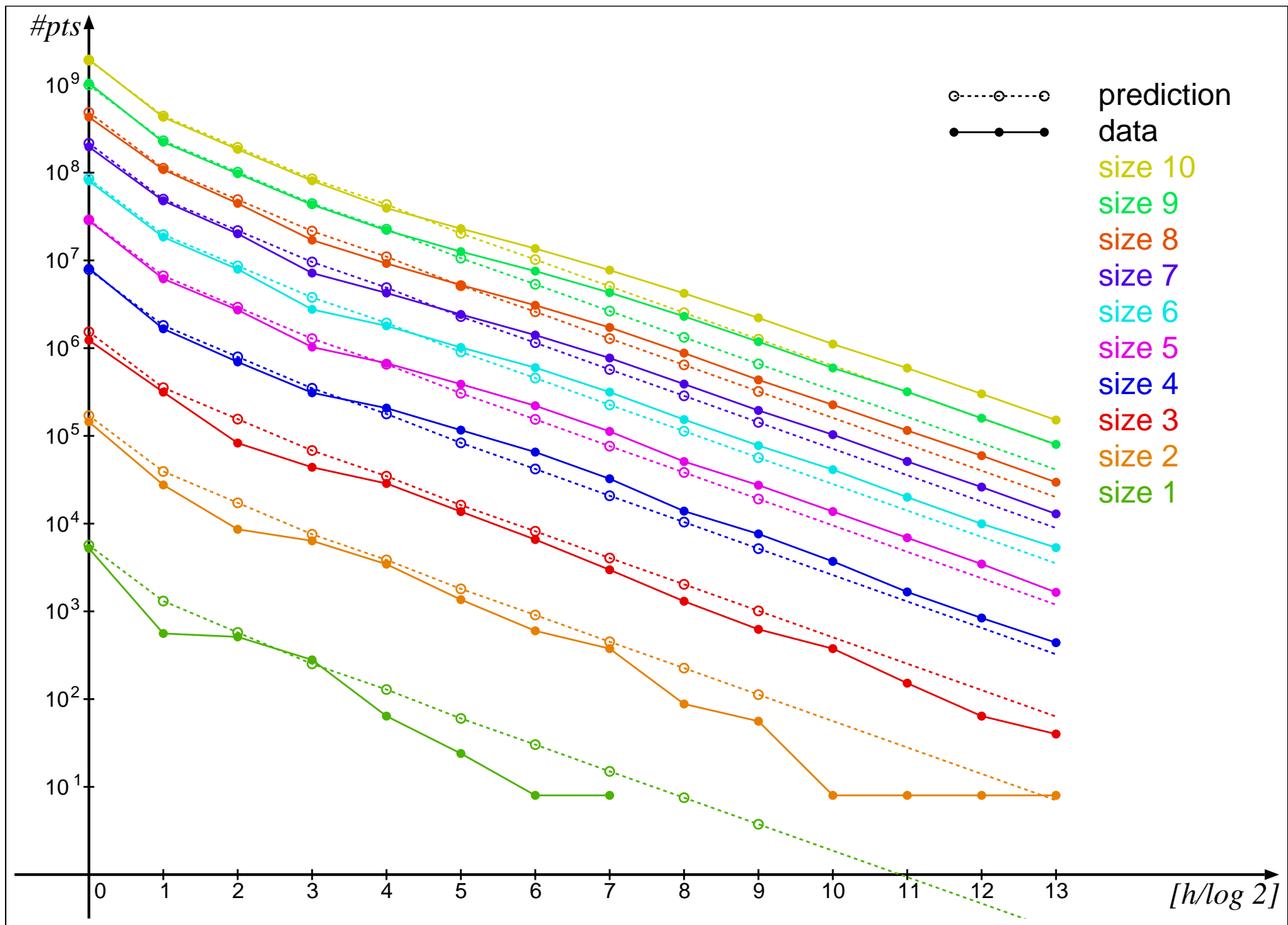
## Some Records (2)

Maximal **Height** of Points ( $H(P) \leq 10^{100}$ )

size of curves	$N = 1$	$N = 2$	$N = 3$
max. $H(P)$	145	10711	209040
max. $H(P)/N^{13/2}$	145.00	118.34	165.55

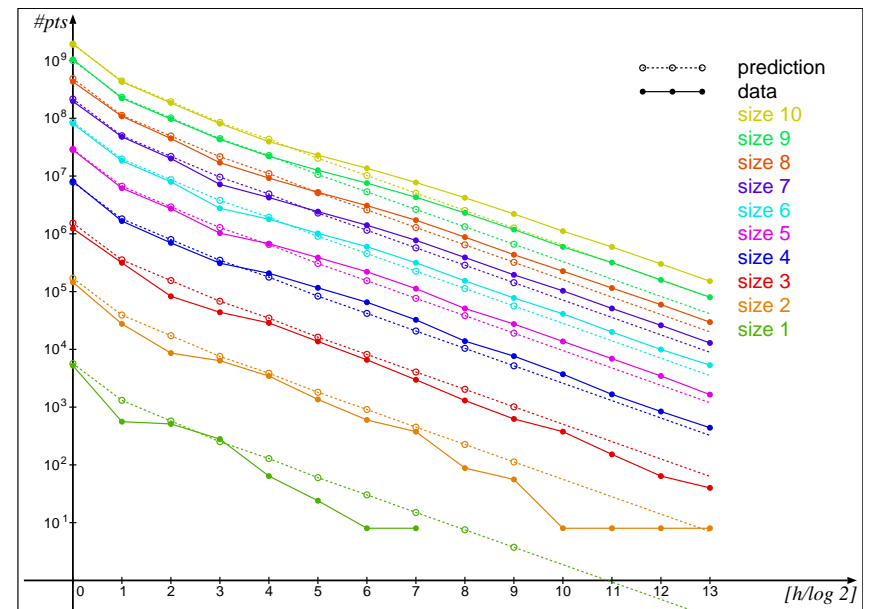
The record point on  $y^2 = x^6 - 3x^4 - x^3 + 3x^2 + 3$

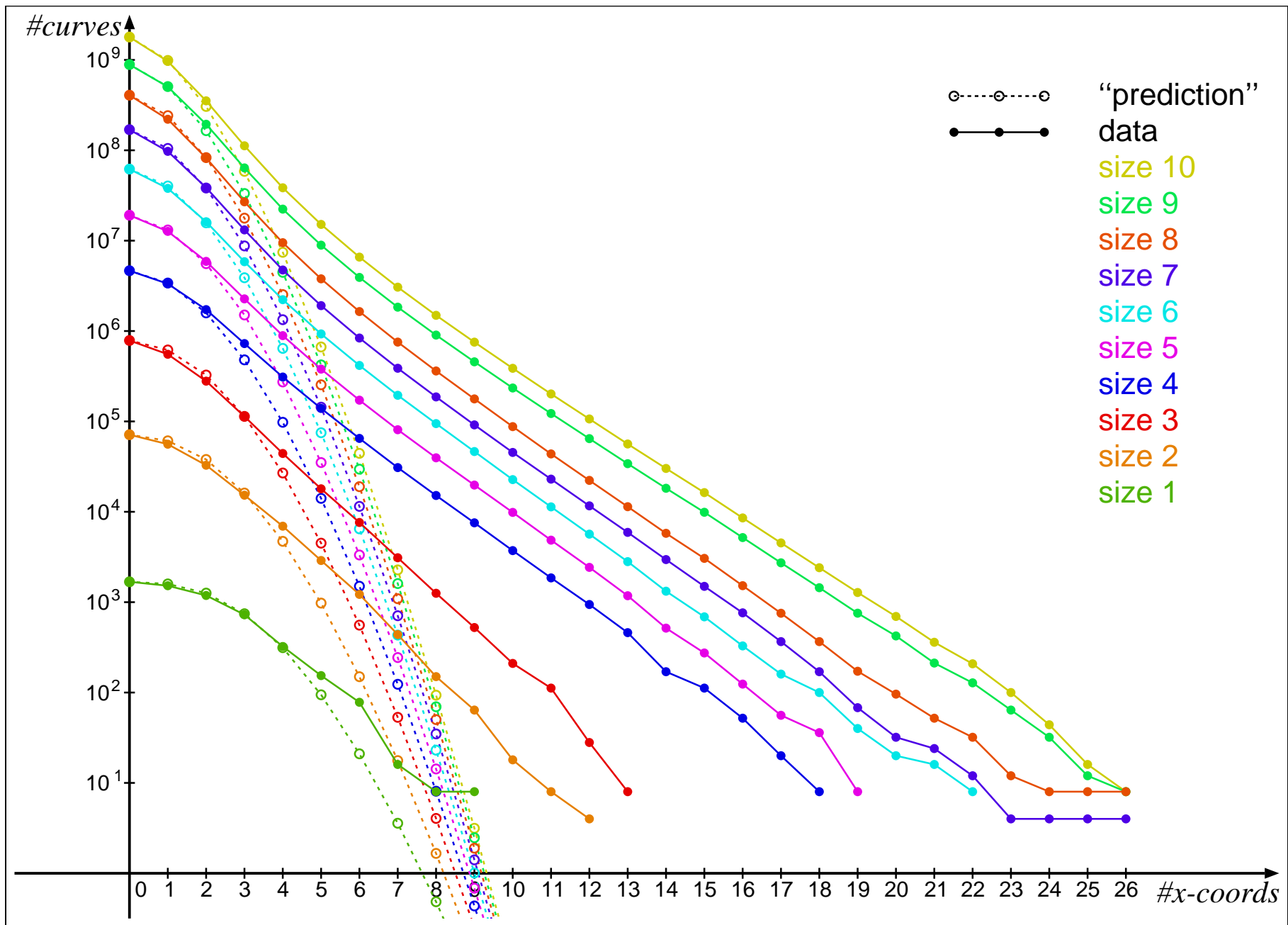
has  $x = -\frac{58189}{209040}$



# Height Distribution of Points

- We count points  $P$  with  $2^n \leq H(P) < 2^{n+1}$ .
- We expect this number to be proportional to  $2^{-n}$ .
- We observe a reasonably **good fit**, but for  $N \geq 4$ , there are “too many” large points. **Explanation?**





## Number of $x$ -Coordinates

- We count curves  $C$  with  $\#\pi(C(\mathbb{Q})) \geq k$ .
- We expect this number to be proportional to  $N^{-k/2}$  (at least for  $k \leq 7$ ).
- We observe that the numbers **drop much more slowly** than expected: Given that there are  $n$  point pairs already, there is a  $\approx 50\%$  chance for another pair.
- This **needs to be investigated!**

# A Pre-Conjecture

Generalizing to larger  $N$ , this leads to the following

**Expectation.**  $\max\{\#\pi(C(\mathbb{Q}))\} \approx c \log(2N + 1).$

Here is a table for  $N \leq 10$ .

size of curves $N$	1	2	3	4	5	6	7	8	9	10
max. $\frac{\#\pi(C(\mathbb{Q}))}{\log(2N+1)}$	8.2	7.5	6.7	8.2	7.9	8.6	9.6	9.2	8.8	8.6

A few **examples** of curves with many points:

Author	$N$	$\#\pi(C(\mathbb{Q}))$	$\frac{\#\pi(C(\mathbb{Q}))}{\log(2N+1)}$
Stahlke	282	63	9.94
Stahlke	249094440	183	9.14
Keller and Kulesz	22999624761	294	11.97

# A Couple of Remarks

## Remark 1.

By work of Caporaso, Harris, and Mazur,

Lang's conjecture on rational points on varieties of general type implies that  $\#C(\mathbb{Q})$  is **uniformly bounded** for curves  $C/\mathbb{Q}$  of fixed genus.

## Remark 2.

One can easily construct families of genus 2 curves with

- at least 14 point pairs and no extra automorphisms,
- at least 16 point pairs and an extra involution,
- at least 24 point pairs and a large automorphism group,

whereas the independence assumption for many points would predict that there are only **finitely many** curves in total with more than 14 point pairs.



# Summary

- For “small” curves  $C$ , we can decide if  $C(\mathbb{Q})$  is empty or not, and we can find all points up to very large height.
- We can use heuristic considerations that lead to expectations regarding the number and size of rational points.
- In a massive computation, we found all rational points of height  $\leq 16383$  on all curves of size  $\leq 10$ .
- The experimental data confirm some of the expectations, but are in disagreement with others.
- It is an interesting challenge to explain the discrepancies!