

# Cassels-Tate pairing on 2-Selmer groups of elliptic curves

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# Outline

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- 2 Higher descents and the Cassels-Tate pairing
- 3 Definition I
- 4 Effectively computing CTP
- 5 Definition II
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## Some notations

- Fix a number field  $k$ .
- $f := X^3 + aX + b = (X - e_1)(X - e_2)(X - e_3) \in k[X]$ , with  $e_1 \neq e_2 \neq e_3 \neq e_1 \in \bar{k}$ .

- Define

$$E := Y^2 = f(X). \quad (1)$$

- For a place  $v$  of  $k$ ,  $k_v$  denotes its completion with respect to  $v$  and  $\mathfrak{k}_v$  the residue field.
- Let  $E_v$  denote the curve  $E$  defined over  $k_v$ .
- $G_F$  we will denote the absolute galois group of the field  $F$ .

## Some theoretical aspects

### Theorem (Mordell-Weil)

$$E(k) \cong E(k)_{tors} \oplus Z^{r_E},$$

where  $\#E(k)_{tors} < \infty$

### Theorem (Lutz-Nagell)

$P := (x, y) \in E(\mathbb{Q})_{tors}$ , then  $x, y \in \mathbb{Z}$ , and  $y^2 | 4a^3 - 27b^2$  or  $y = 0$ .

- A faster way to compute  $E(k)_{tors}$  is to use the injection  $E(k)_{tors} \hookrightarrow \bar{E}_v(\mathfrak{k}_v)$ ,  $\bar{E}_v$  denotes  $E$  over  $\mathfrak{k}_v$  for a place of good reduction.
- No unconditional algorithm to compute  $r_E$  is known.
- Assuming BSD,  $r_E$  may be computed using  $\text{ord}(L(E, s = 1))$ .

## Upper and lower bounds

- **Lower bound:** Find points and check for independence.
- **Upper bound:** Use descent.
- See if they match!

The Kummer sequence:

$$0 \longrightarrow E[n] \hookrightarrow E \longrightarrow E \longrightarrow 0 \quad (2)$$

Applying galois cohomology:

$$\begin{array}{ccccccc}
 0 & \longrightarrow & \frac{E(k)}{nE(k)} & \hookrightarrow & H^1(G_k, E[n]) & \longrightarrow & H^1(G_k, E)[n] \longrightarrow 0 \\
 & & \downarrow & & \downarrow & \searrow \alpha & \downarrow \\
 0 & \longrightarrow & \prod_v \frac{E(k_v)}{nE(k_v)} & \hookrightarrow & \prod_v H^1(G_{k_v}, E[n]) & \longrightarrow & \prod_v H^1(G_{k_v}, E_v)[n] \longrightarrow 0,
 \end{array}$$

## Descent sequence

We have the  $n$ -descent sequence.

$$0 \longrightarrow \frac{E(k)}{nE(k)} \hookrightarrow \text{Sel}^{(n)}(E) \longrightarrow \text{III}[n] \longrightarrow 0, \quad (3)$$

where  $n^{\text{th}}$  Selmer group

$$\text{Sel}^{(n)}(E) = \ker(\alpha),$$

and the Tate-Shafarevich group

$$\text{III} := \ker \left( H^1(G_k, E) \longrightarrow \prod_v H^1(G_{k_v}, E_v) \right).$$

- $\text{Sel}^{(n)}(E)$  is finite and effectively computable in principle.
- $\text{rank}_{\mathbb{F}_p}(\text{Sel}^{(p)}(E))$  bounds the  $\text{rank}_{\mathbb{F}_p}(\frac{E(k)}{pE(k)})$ .
- If  $\text{III}[n]$  is trivial then  $\text{Sel}^{(n)}(E) \cong \frac{E(k)}{nE(k)}$ .

# Higher descent

- If  $(m, n) = 1$ ,  $\text{Sel}^{(mn)}(E) \cong \text{Sel}^{(n)}(E) \times \text{Sel}^{(m)}(E)$ .
- If  $l \geq 2$ , then  $\text{Sel}^{p^l}(E)$  is known as **higher  $p$ -descent**.
- We have the following commutative diagram:

$$\begin{array}{ccc}
 \frac{E(k)}{p^l E(k)} & \hookrightarrow & \text{Sel}^{(p^l)}(E) \\
 \downarrow & & \downarrow p \\
 \frac{E(k)}{p^{l-1} E(k)} & \hookrightarrow & \text{Sel}^{(p^{l-1})}(E)
 \end{array} \tag{4}$$

We have

$$\frac{E(k)}{pE(k)} \subseteq p\text{Sel}^{(p^2)}(E) \subseteq \text{Sel}^{(p)}(E). \tag{5}$$

## Cassels-Tate pairing (CTP)

- CTP (denoted by  $\langle \cdot, \cdot \rangle_{CT}$ ) was defined by **J.W.S. Cassels** for elliptic curves.
- **John Tate** generalized it to abelian varieties.
- CTP has following properties:
  - $\langle \cdot, \cdot \rangle_{CT} : \text{III} \times \text{III} \longrightarrow \mathbb{Q}/\mathbb{Z}$ .
  - CTP is an anti-symmetric pairing.
  - For  $\alpha \in \text{III}[n]$ ,  $\langle \beta, \alpha \rangle_{CT} = 0 \iff \beta \in n\text{III}$ .

Pulling back  $\langle \cdot, \cdot \rangle_{CT}$  on  $\text{Sel}^{(n)}(E)$ , we have:

### Theorem (Cassels)

$\langle \alpha, \beta \rangle_{CT} = 0$  for all  $\beta \in \text{Sel}^{(n)}(E)$ ,  $\iff \alpha \in n\text{Sel}^{(n^2)}(E)$ .



## Weil-pairing based definition

- Consider the “evaluation” based pairings:
  - $\langle \cdot, \cdot \rangle_1 : (\text{Princ}(E) \times \text{Div}^0(E))^\perp \longrightarrow \overline{k}^\times$ ,
  - $\langle \cdot, \cdot \rangle_2 : (\text{Div}^0(E) \times \text{Princ}(E))^\perp \longrightarrow \overline{k}^\times$ ,
 with  $\langle \text{div}(f), D \rangle_1 = \langle D, \text{div}(f) \rangle_2 = \prod_P f(P)^{v_P(D)}$ .
- $\langle \cdot, \cdot \rangle_1 = \langle \cdot, \cdot \rangle_2$  on  $(\text{Princ}(E) \times \text{Princ}(E))^\perp$  (Weil-reciprocity!).
- Let  $e_n(\cdot, \cdot) : E[n] \times E[n] \longrightarrow \mu_n$ , denote the **Weil-pairing** with

$$e_n(P, Q) = \langle np, q \rangle_1 - \langle p, nq \rangle_2,$$

where  $p, q$  denote the representatives of  $P, Q$  (resp.) in  $\text{Div}^0(E)$ .

- Let  $\cup$  denote the cup product on  $H^*(G_k, E[n])$  induced by Weil-pairing.

## Weil-pairing based definition (cont.)

Let  $\alpha, \alpha' \in \text{Sel}^{(n)}(E)$ .

- **Global part:**

- Lift  $\alpha, \alpha'$  to  $\mathfrak{a}, \mathfrak{a}' \in Z^1(G_k, E[n])$ .
- Let  $a \in C^1(G_k, E[n^2])$ , with  $na = \mathfrak{a}$ .
- $\partial a \cup \mathfrak{a}' \in Z^3(G_k, \bar{k}^\times) \implies \partial a \cup \mathfrak{a}' = \partial \epsilon$ , for  $\epsilon \in C^2(G_k, \bar{k}^\times)$ .
- The above statements follow from galois cohomology on

$$0 \longrightarrow E[n] \hookrightarrow E[n^2] \longrightarrow E[n] \longrightarrow 0$$

and that  $H^3(G_k, \bar{k}^\times)$  is trivial.

- **Local part:**

- $\alpha_v = 0 \implies \exists P_v \in E_v$ , with  $\partial P_v = \alpha_v$ .
- Choose  $Q_v \in E_v$ , with  $nQ_v = P_v$  and let  $q_v := \partial Q_v$ .
- $a_v - q_v$  take values in  $E[n]$ .
- Define  $\gamma_v := (a_v - q_v) \cup \mathfrak{a}'_v - \epsilon_v \in Z^2(G_{k_v}, \bar{k}_v^\times)$ .

## Weil-pairing based definition (contd.)

$\gamma_v$  represents some class  $c_v \in H^2(G_{k_v}, \bar{k}_v^\times) \cong \text{Br}(k_v)$ .

### Definition

For  $(\alpha, \alpha') \in \text{Sel}^{(n)}(E) \times \text{Sel}^{(n)}(E)$  we have:

$$\langle \alpha, \alpha' \rangle_{CT} = \sum_v \text{inv}_v(c_v)$$

**Effectiveness of CTP:** Cassels effectively defined a pairing  $\langle \cdot, \cdot \rangle_{Cas}$  on  $\text{Sel}^{(2)}(E) \times \text{Sel}^{(2)}(E)$ , with properties same as CTP.

## Known results

- **Fischer, Schaefer and Stoll**, showed that  $\langle \cdot, \cdot \rangle_{Cas} = \langle \cdot, \cdot \rangle_{CT}$  on  $\text{Sel}^{(2)}(E) \times \text{Sel}^{(2)}(E)$ .
- **Swinnerton-Dyer** extended this approach to compute CTP between  $\text{Sel}^{(2)}(E)$  and  $\text{Sel}^{(2^n)}(E)$ .
- **Fischer and Newton** computed CTP on  $\text{Sel}^{(3)}(E)$ .
- **van Beek and Fischer** compute CTP on Selmer groups of odd prime degree isogeny.
- The above computations were based on Weil-pairing based definition.

## Known results (contd.)

- **Fischer and Donnelly** used homogenous space based definition to compute CTP on  $\text{Sel}^{(2)}(E)$ .
- **Tom Fischer** has a similar approach to compute CTP on  $\text{Sel}^{(3)}(E)$  using homogenous space definition.
- CTP can be defined in general for  $\text{III}(A) \times \text{III}(A^\vee)$ , for an abelian variety  $A$  and its dual  $A^\vee$ .
- We compute CTP using **Albanese-Albanese** definition given by **Poonen and Stoll**.
- We aim to compute CTP on jacobians of genus 2 curves.

## Albanese-Albanese definition

Let  $\alpha, \alpha' \in \text{Sel}^{(n)}(E)$ .

- **Global part:**

- Lift  $\alpha, \alpha'$  to  $\mathfrak{a}, \mathfrak{a}' \in C^1(G_k, \text{Div}^0(E))$ .
- $\partial\mathfrak{a}, \partial\mathfrak{a}'$  take values in  $\text{Princ}(E)$ .
- Let  $\eta := \langle \partial\mathfrak{a}, \mathfrak{a}' \rangle_1 - \langle \mathfrak{a}, \partial\mathfrak{a}' \rangle_2 \in Z^3(G_k, \bar{k}^\times) \implies \eta = \partial\epsilon$ , for  $\epsilon \in C^2(G_k, \bar{k}^\times)$ .
- The above statements follow using the galois cohomology on Kummer sequence and on

$$0 \longrightarrow \text{Princ}(E) \hookrightarrow \text{Div}^0(E) \longrightarrow \text{Pic}^0(E) \longrightarrow 0,$$

and using  $H^3(G_k, \bar{k}^\times)$  is trivial.

- **Local part:**

- There exists  $P_v \in E_v$ , with  $\partial P_v = \alpha_v$ .
- Lift  $P_v$  to a degree zero divisor  $\mathfrak{p}_v$ , and  $\mathfrak{a}_v - \partial\mathfrak{p}_v$  takes values in  $\text{Princ}(E)$ .
- Consider  $\gamma_v := \langle (\mathfrak{a}_v - \mathfrak{p}_v), \mathfrak{a}'_v \rangle_1 - \langle \mathfrak{p}_v, \partial\mathfrak{a}'_v \rangle_2 - \epsilon_v$ .

## Albanese-Albanese definition (contd.)

$\gamma_v$  represents some class  $c_v \in H^2(G_{k_v}, \bar{k}_v^\times) \cong \text{Br}(k_v)$ .

### Definition

For  $(\alpha, \alpha') \in \text{Sel}^{(n)}(E) \times \text{Sel}^{(n)}(E)$  we have:

$$\langle \alpha, \alpha' \rangle_{CT} = \sum_v \text{inv}_v(c_v)$$

### Theorem (Poonen and Stoll)

*Weil-pairing based definition and Albanese-Albanese definition of the Cassels-Tate pairing are equal.*

We prove the following theorem:

### Theorem

Let  $\alpha, \alpha' \in \text{Sel}^{(2)}(E)$ , represented by  $(\beta_1, \beta_2, \beta_3)$ ,  $(\beta'_1, \beta'_2, \beta'_3)$ , with  $\beta_1\beta_2\beta_3 \in k^2$ , and  $\beta'_1\beta'_2\beta'_3 \in k^2$  and  $\beta_i, \beta'_i \in k(\mathbf{e}_i)$ .

$(-1)^{2\langle \alpha, \alpha' \rangle_{CT}} = \prod_v [\alpha, \alpha']_v$ , where

$$[\alpha, \alpha']_v = \begin{cases} \prod_{i=1}^3 (\delta_{v,i}, \beta'_i)_{k_v}, & f \text{ splits over } k, \\ (\delta_{v,1}, \beta'_1)_{k_v} (\delta_{v,2}, \beta'_2)_{k_v(\mathbf{e}_2)} & \mathbf{e}_1 \in k \text{ and } [k(\mathbf{e}_2) : k] = 2, \\ (\delta_{v,1}, \beta'_1)_{k_v(\mathbf{e}_1)} & [k(\mathbf{e}_1) : k] \geq 3, \end{cases}$$

where  $\delta_{vi} \in k_v(\mathbf{e}_i)$ .

### Corollary

The above pairing exactly matches the  $\langle \cdot, \cdot \rangle_{Cas}$ .



# A bottleneck and exploiting freedom of choices!

- An important step is to find  $\epsilon \in C^2(G_k, \bar{k}^\times)$  such that  $\partial\epsilon = \eta$ .
- Solving "skewed" linear equations:

$$\frac{\sigma\epsilon(\tau, \rho)\epsilon(\sigma, \tau\rho)}{\epsilon(\sigma\tau, \rho)\epsilon(\sigma, \tau)} = \eta(\sigma, \tau, \rho).$$

- We exploit the everywhere local solubility of 2-coverings to get global solution to certain norm equations.

**A simplification:** Consider a different lift of  $\alpha'$  to  $\mathfrak{a}$  with values in  $\text{Div}^0(E)$ .

## Exploiting freedom of choices

Assume  $f$  splits over  $k$ .

$$\alpha' := \sum_{i=1}^3 \alpha'_i,$$

with  $\alpha'_i$  being 1-cocycles.

This splits

$$\eta = \sum_{i=1}^3 \eta_i,$$

with  $\eta_i \in H^3(G_k, \bar{k}^\times)$  using bilinearity of cup-product.

Find  $\epsilon_i$ , with

$$\partial \epsilon_i = \eta_i.$$

*Thank You!*